

## Ext/Tor :

Goal: For  $R$  modules  $M, N$ , define

$$\text{Ext}_R^i(M, N)$$
$$\text{Tor}_j^R(M, N)$$

w/  $\text{Ext}_R^0(M, N) = \text{Hom}_R(M, N)$ ,  $\text{Tor}_0^R(M, N) = M \otimes_R N$ ,

- $\text{Ext}_\mathbb{Z}^{>1}(M, N) = 0$  so use shorthand  $\text{Ext}(M, N)$  for  $\text{Ext}_\mathbb{Z}^1(M, N)$ , similar for  $\text{Tor}^{\mathbb{Z}}$ .

## Last time:

• An  $R$ -module  $Q$  is injective  $R$ -module if, for any injective map (of  $R$ -modules)  $f: M \rightarrow N$ ,

and any map  $g: M \rightarrow Q$   $\exists$

$$\text{(SES)} \quad 0 \rightarrow M \xrightarrow{f} N$$

"any  $g$  to  $Q$  extends along injectives".

$$\begin{array}{ccc} & & \swarrow \exists h \text{ with } h \circ f = g \\ & \downarrow & \\ & Q & \end{array}$$

(injective module)

exercise:

$\iff$  if any SES  $0 \rightarrow Q \rightarrow M \rightarrow K \rightarrow 0$  splits.

Remark:  $0 \rightarrow A \xrightarrow{i} B \xrightarrow{j} C \rightarrow 0$  splits iff  $B \cong A \oplus C$  w/  $i$  inclusion,  $j$  projection

$$\iff \exists 0 \rightarrow A \xrightarrow{\underbrace{s}} B \xrightarrow{\underbrace{j \circ s^{-1}}} C \rightarrow 0 \text{ w/ } j \circ s = \text{id}_C \iff$$

$$\exists 0 \rightarrow A \xrightarrow{\underbrace{s'}} B \xrightarrow{\underbrace{j \circ s'^{-1}}} C \rightarrow 0 \text{ w/ } j \circ s' = \text{id}_C$$

exercise

$\iff$  if  $\text{Hom}_R(-, Q)$  is exact.

• An  $R$ -module  $P$  is projective if for any surjection  $f: N \rightarrow M$  and any map  $g: P \rightarrow M$ ,  $\exists$ :

$$\begin{array}{ccc} N & \xrightarrow{f} & M \rightarrow 0 \quad \text{(SES)} \\ & \nwarrow \exists h & \uparrow g \\ & & P \end{array}$$

with  $f \circ h = g$  (projective)

"any  $g$  from  $P$  lifts along surjections."

$\Leftrightarrow$  any <sup>SES</sup>  $0 \rightarrow A \rightarrow B \rightarrow P \rightarrow 0$  is split.

$\Leftrightarrow \text{Hom}(P, -)$  is exact.

Thm: (exercise or look up in a book): For a  $\mathbb{Z}$ -module  $M$  (i.e., an abelian group) (or more generally Mod  $\leftarrow \text{PID}$   $\leftarrow R$ ).

•  $M$  is injective iff it is divisible.

(an abelian group  $G$  is divisible if for any  $g \in G$  and any  $n \in \mathbb{N}$ ,  $g = n \cdot g'$  for some  $g' \in G$ )

ex:  $\mathbb{Q}$ , non-ex:  $\mathbb{Z}$ , or  $\mathbb{Z}/2$

$\leftarrow$  but  $\mathbb{Z}/2$  is injective as a  $\mathbb{Z}/2$ -module!

•  $M$  is projective iff it is free. (for any  $R$ , free  $\Rightarrow$  projective, not always  $\Leftarrow$ )

Cor: For a projective  $\mathbb{Z}$ -module  $P$ , given injection  $0 \rightarrow P' \rightarrow P$  (i.e.,  $\hookrightarrow$  a submodule),  $P'$  is projective too. (subgroups of free abelian groups are free abelian)

Similarly, if  $Q$  injective  $\mathbb{Z}$ -module,  $Q \rightarrow Q' \rightarrow 0$ ,  $Q'$  injective too.

A projective resolution of an  $R$ -module  $M$  is an exact sequence

$$\rightarrow \dots \rightarrow P_2 \rightarrow P_1 \rightarrow M \rightarrow 0$$

resp.  
(injective resolution:

$$0 \rightarrow Q \rightarrow I_1 \rightarrow I_2 \rightarrow \dots)$$

If turns out those resolutions always exist, and

as a consequence of Thm/Cor above, we deduce:

Any  $\mathbb{Z}$ -module (i.e., abelian group)  $M$  has a two-step projective (resp. injective) resolution

$$0 \rightarrow P_1 \rightarrow P_0 \xrightarrow{f} M \rightarrow 0. \text{ (why? Pick surjection } f: P_0 \rightarrow M \text{ } \& \text{ have}$$

$$0 \rightarrow \ker f \rightarrow P_0 \rightarrow M \rightarrow 0$$

$$\text{(resp. } 0 \rightarrow M \rightarrow I_0 \rightarrow I_1 \rightarrow 0 \text{).}$$

?  
arg. by cor)

Def: Given  $R$ -modules  $M, N$ , pick projective resolution of  $M$

$$\dots \xrightarrow{f_1} P_1 \xrightarrow{f_0} P_0 \rightarrow M \rightarrow 0,$$

take  $\text{Hom}_R(-, N)$  of all  $P_i$ 's &  $f_i$ 's between them (but not of  $M$ ):

$$\dots \xleftarrow{f_1^*} \text{Hom}(P_1, N) \xleftarrow{f_0^*} \text{Hom}(P_0, N) \leftarrow 0$$

This is no longer exact; however it's a chain complex:  $f_{i+1}^* \circ f_i^* = 0$ . (so  $\ker f_{i+1}^* \supseteq \text{im } f_i^*$ ).

We define  $\text{Ext}_R^i(M, N) := \frac{\ker f_i^*}{\text{im } f_{i-1}^*}$ . Note  $\text{Ext}_R^0(M, N) = \ker f_0^* \cong \text{Hom}_R(M, N)$

(b/c  $P_1 \xrightarrow{f_0} P_0 \rightarrow M \rightarrow 0$  induces an exact sequence

$$0 \rightarrow \text{Hom}(M, N) \rightarrow \text{Hom}(P_0, N) \xrightarrow{f_0^*} \text{Hom}(P_1, N)$$

This a priori depends on a choice (of proj. resolution);

however

Thm:  $\text{Ext}_R^i(M, N)$  are independent of choices made and functorial in  $M/N$ .

Remark: Can also define  $\text{Ext}_R^i(M, N)$  by injectively resolving  $N$ , exercise that the result is unchanged.

Cor: For  $\mathbb{Z}$ -modules  $M, N$ ,  $\text{Ext}_{\mathbb{Z}}^i(M, N) = 0$  for  $i > 1$ .

(b/c  $M$  admits a two-term proj. resolution  $0 \rightarrow P_1 \rightarrow P_0 \rightarrow M \rightarrow 0$ ).

Similarly, to define  $\text{Tor}_j^R(M, N)$ , projectively resolve either  $M$  or  $N$ :

$(\dots \rightarrow P_1 \rightarrow P_0) \rightarrow M \rightarrow 0$ , then take  $P_i \otimes_R N$ ; again the result is a chain complex (no longer exact);  $j$ th homology is  $\text{Tor}_j^R(M, N)$ .

& same Thm: functorial + independent of choice.

Both these theorems follow from "homotopically unique functoriality of projective resolutions":

Thm: Say  $P_\bullet \rightarrow M$  is a proj. resolution and  $\bar{f}: M \rightarrow N$  is a map (of  $R$ -modules), and  $Q_\bullet \rightarrow N$  a proj. resolution. Then there is a chain map  $f_\bullet: P_\bullet \rightarrow Q_\bullet$  'lifting'  $\bar{f}$ ,

meaning:

$$\begin{array}{ccccccc} \dots & \xrightarrow{\partial^P} & P_1 & \xrightarrow{\partial^P} & P_0 & \rightarrow & M \rightarrow 0 \\ & & \vdots & & \vdots & & \downarrow \bar{f} \\ \dots & \xrightarrow{\partial^Q} & Q_1 & \xrightarrow{\partial^Q} & Q_0 & \rightarrow & N \rightarrow 0 \end{array}$$

making diagram commute.

and  $f$  is unique up to homotopy equivalence (meaning if  $f, f': P \rightarrow Q$  then  $\exists h: P \rightarrow Q$  s.t.  $f - f' = \partial \circ h + h \circ \partial$ )

with  $f - f' = \partial \circ h + h \circ \partial$ .

— This induces, for  $f: M \rightarrow N$ , a map  $f^*$  between cochain complexes computing  $\text{Ext}_R^i(M, S)$  &  $\text{Ext}_R^i(N, S)$ .  
 inde of choices, let

To see from this that e.g.,  $\text{Ext}$  is  $\wedge$   $N=M, \bar{f}=\text{id}, P, Q$  two different resolutions of  $M$ ,

(\*)  $\Rightarrow$  get a chain map  $f: P \rightarrow Q$  lifting  $\bar{f}$ , & hence (by taking  $\text{Hom}_R(-, S)$ )  
 a chain map between Ext complexes,  $f^*$

we also get a map  $g: Q \rightarrow P$  & get  $g^*$  other way between Ext complexes:

By homotopical uniqueness,  $f \circ g$  &  $g \circ f$  are each homotopic to  $\text{id}_Q, \text{id}_P$  respectively,

hence  $f^*, g^*$  induce isomorphisms on Ext groups.

Sketch of proof of theorem: idea is to inductively use the projectivity condition:

inductively say we've constructed  $f_i, \dots, f_0$ . (base case:  $\bar{f}$  itself  $i = -1$ )

$$\begin{array}{ccccccc} \dots & \rightarrow & P_{i+1} & \xrightarrow{\partial^P} & P_i & \xrightarrow{\partial^P} & \dots \rightarrow P_0 \rightarrow M \\ & & \downarrow f_{i+1} & & \downarrow f_i & & \downarrow f_0 & \downarrow \bar{f} \\ \dots & \rightarrow & Q_{i+1} & \xrightarrow{\partial^Q} & Q_i & \xrightarrow{\partial^Q} & \dots \rightarrow Q_0 \rightarrow N \end{array}$$

want:  $\partial^Q \circ f_{i+1} = f_i \circ \partial^P$ . Since sequences at top and bottom are exact, can restrict to:

$$\begin{array}{ccc} P_{i+1} \rightarrow \ker(\partial_i^P) \rightarrow 0 & & 0 \rightarrow \ker(\partial_i^P) \rightarrow P_i \rightarrow \dots \rightarrow 0 \\ \downarrow f_{i+1} & \text{map from proj.} & \downarrow f_i & & \downarrow f_i & & \downarrow f_i & & \downarrow f_{i-1} \\ Q_{i+1} \rightarrow \ker(\partial_i^Q) \rightarrow 0 & & & \text{and} & & & 0 \rightarrow \ker(\partial_i^Q) \rightarrow Q_i \rightarrow \dots \rightarrow 0 \end{array}$$

*surjection*

Uniqueness up to homotopy? exercise. similar inductive level-by-level construction of

$$\begin{array}{ccccccc} \dots & \rightarrow & P_{i+1} & \rightarrow & P_i & \rightarrow & \dots \\ & & \downarrow f_{i+1} - f'_{i+1} & & \downarrow f_i - f'_i & & \\ \dots & \rightarrow & Q_{i+1} & \rightarrow & Q_i & \rightarrow & \dots \end{array}$$

satisfying  $\partial h + h \partial = f - f'$ .