

Day 2 Talk 4: Nick S, Combinatorial Feynman Categories

Setup: $(M, \mathcal{I}, \omega, \Theta)$ $\rightarrow \omega = d\Theta$.
Riemann surface (with boundary) — any area form

Defn: Let (L_1, \dots, L_n) be a collection of exact Lagrangian S^1 's embedded in M .

(Don't have to worry about bubbling, can set $\epsilon > 1$ in Λ).

Exactness $\Rightarrow L_i$ not nullhomologous.

$\mathcal{F}(L_1, \dots, L_n) \rightarrow$ over field \mathbb{K} , $\text{char } \mathbb{K} = 2$.

(let's not worry about gradings right now).

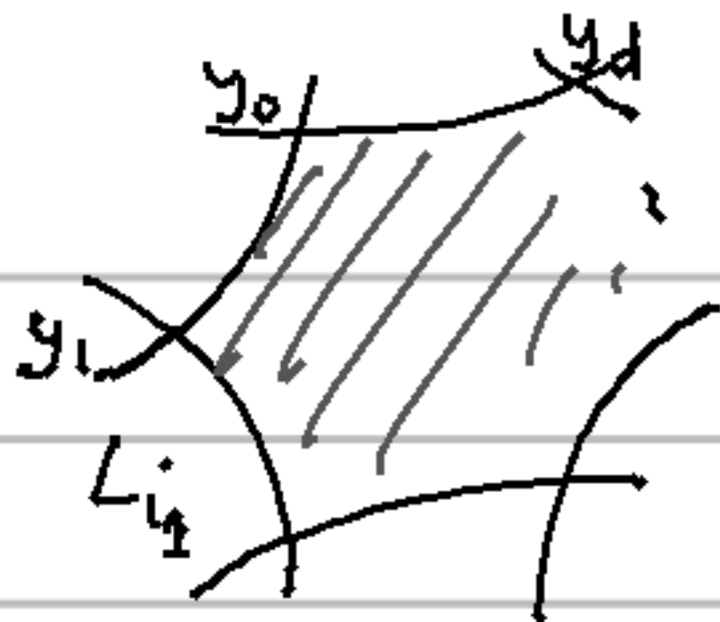
$$\text{Ob } \mathcal{F} = \{L_i\}$$

$$\text{hom}(L_i, L_j) = \begin{cases} \mathbb{K}^{L_i \cap L_j} & i < j \\ \mathbb{K} \cdot e & i = j \\ 0 & i > j \end{cases}$$

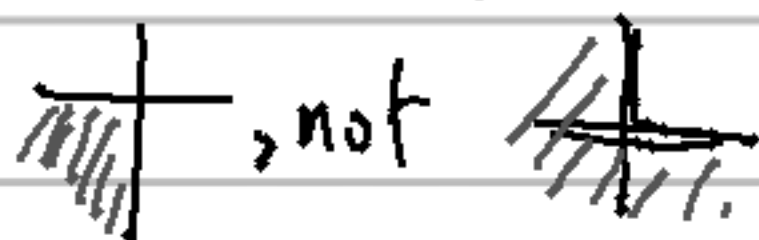
Define η^d : Given $y_0 \in L_{i_0} \cap L_{i_1}$

$$y_k \in L_{i_{k-1}} \cap L_{i_k} \quad (k = 1, \dots, d)$$

Let $n_d(y_0, \dots, y_d) = \#$ immersed d -gons in M like this:



we want
convex corners



Example of immersed 2-gon:



$$y^d: \text{hom}(L_{i_d}, L_{i_d}) \otimes \dots \otimes \text{hom}(L_{i_0}, L_{i_1})$$

$$\downarrow$$

$$\text{hom}(L_{i_0}, L_{i_d})$$

$$y^d(y_d \otimes \dots \otimes y_1) = \sum_{y_0} n_d(y_0, \dots, y_1) y_0$$

(= 0 unless $i_0 < \dots < i_d$)

(why is this the same as $\mathcal{F}^{\rightarrow}(L_1, \dots, L_n)$?)

Ans: essentially Riemann mapping thm.)

This is an \mathcal{A}_{y_0} category: (using codomain picture).

non-commutative non-unital co-algebra

$$T(L_1, \dots, L_n) = \bigoplus_{d \geq 1} T_d$$

$$T_d = \bigoplus_{i_0 < \dots < i_d} \text{hom}(L_{i_0}, L_{i_1}) \otimes \dots \otimes \text{hom}(L_{i_{d-1}}, L_{i_d})^*$$

(Define $\alpha \otimes \beta = 0$ if wrong ordering.)

T_d has a basis

$$\left\{ a_1 \otimes \dots \otimes a_d \mid a_j \in L_{i_j} \wedge L_{i_{j+1}}, \right. \\ \left. i_1 < \dots < i_d \right\}$$

(really this is a dual basis).

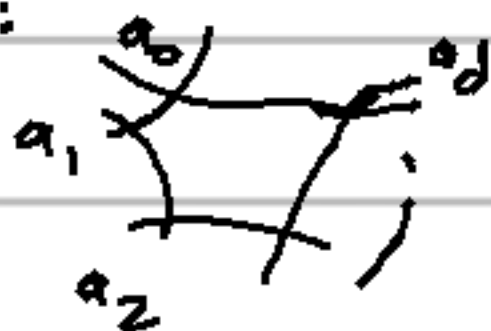
Define $\delta: T \rightarrow T$ by

$$\delta(a_0) = \sum_j u_j(a_0, \dots, a_d) a_1 \dots a_d$$

$$(\text{=} u_1^* + u_2^* + \dots)$$

For each polygon like

this:



get a contribution to $\delta(a_0)$.

Extend $\delta|_{T_1}$ to T by
Leibniz rule.

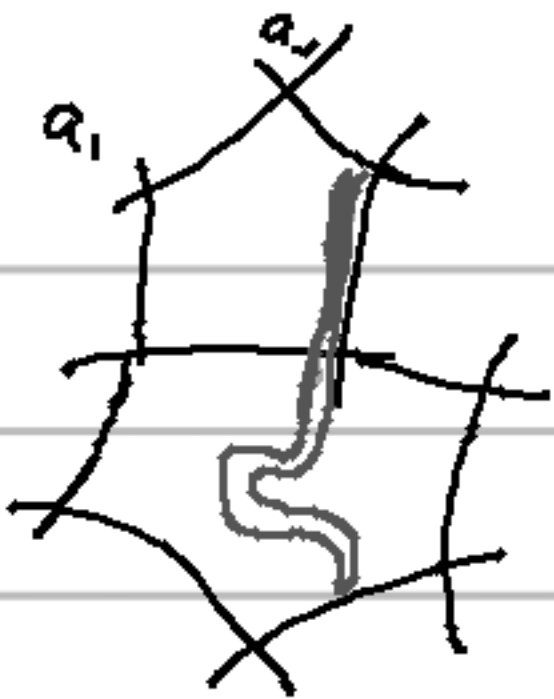
Prop: $\delta^2 = 0$

$$\delta^2 a_0 = \delta \sum_{\text{polys}} a_1 \dots a_d$$

$$= \sum_i \sum_{\text{polys}} a_1 \dots \delta a_i \dots a_d$$

$$= \sum_{\text{polys}} \sum_{\text{polys}} a_1 \dots a_{i-1} b_1 \dots b_p a_{i+1} \dots a_d$$

i.e. sum of things that look like



sum of polygons w/
 one convex corner
 i.e. I don't finish,
 so fake boundary
 & follow out.

\Rightarrow terms in $\delta^2 a_0$ cancel in pairs

$\Rightarrow \delta^2 a_0 = 0 \Rightarrow \delta^2 = 0$ (two things that

Cor: μ^d define A_{∞} structure.

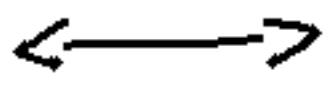
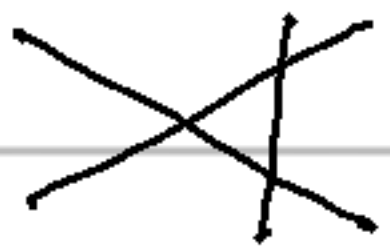
(exact Lag's key to ensuring a finite sum).

go into the
 concave vertex
 must come
 out).

Prop: Hamilton isotopy of L_i changes $\mathcal{F}(L_1, \dots, L_n) \rightarrow$

by a quasi-isomorphism.

PK:
 I.



Moves we worry
 about

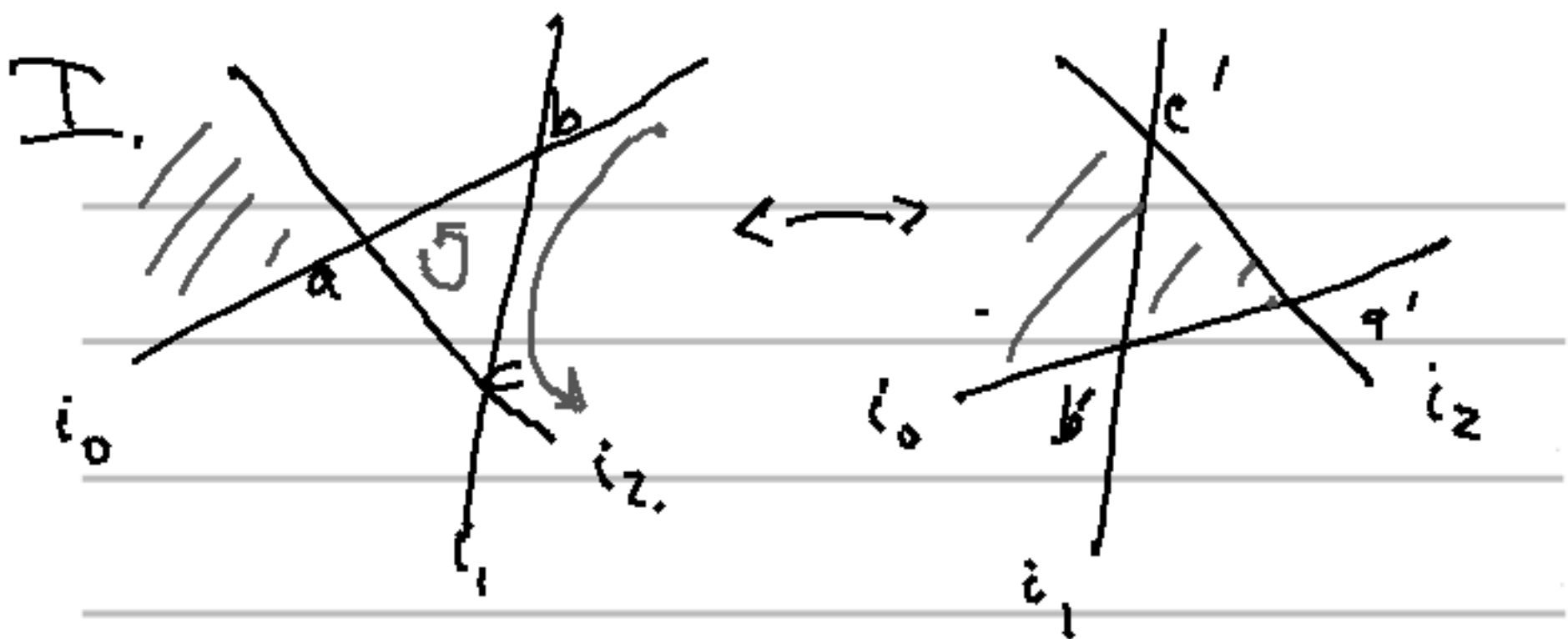
II.



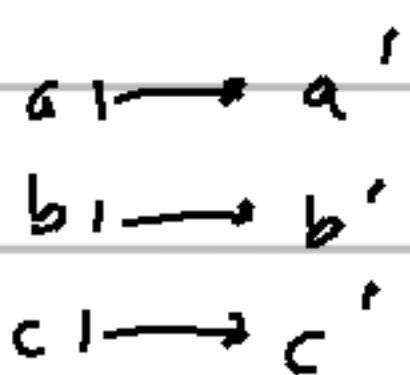
(i.e. Reidemeister moves 2 & 3 but don't have

I b/c Lag's are embedded).

Let's look at one of these cases.



WLOG i_0 is smallest,
 if not ($i_0 < i_1 < i_2$),



if $i_0 < i_1 < i_2$:

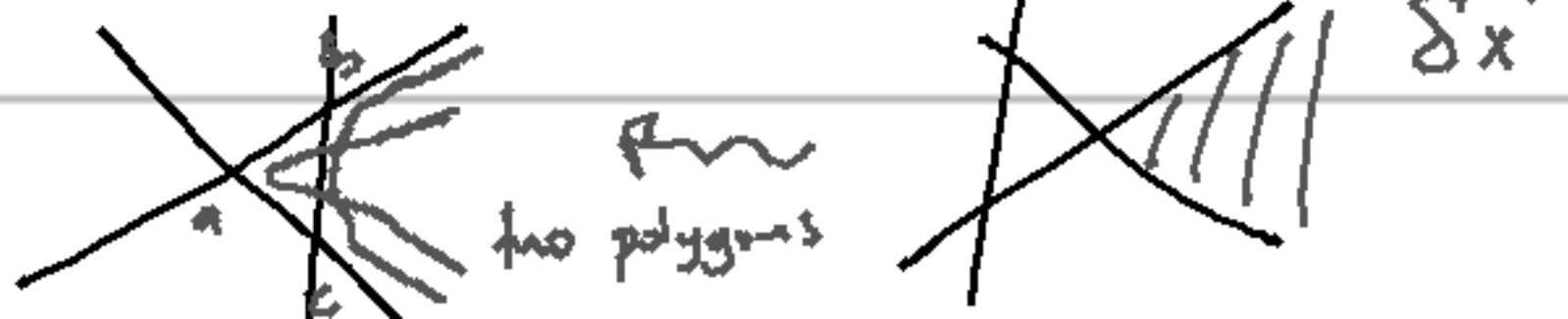
$$f: a \mapsto a' + b'c'$$

$$x \mapsto x' \text{ for all other } x.$$

$$a + bc \longleftarrow a': g$$

$$x \longleftarrow x'$$

f, g are chain maps: $x \neq a$,
 $\delta' f x = \delta' x' = f \delta x$



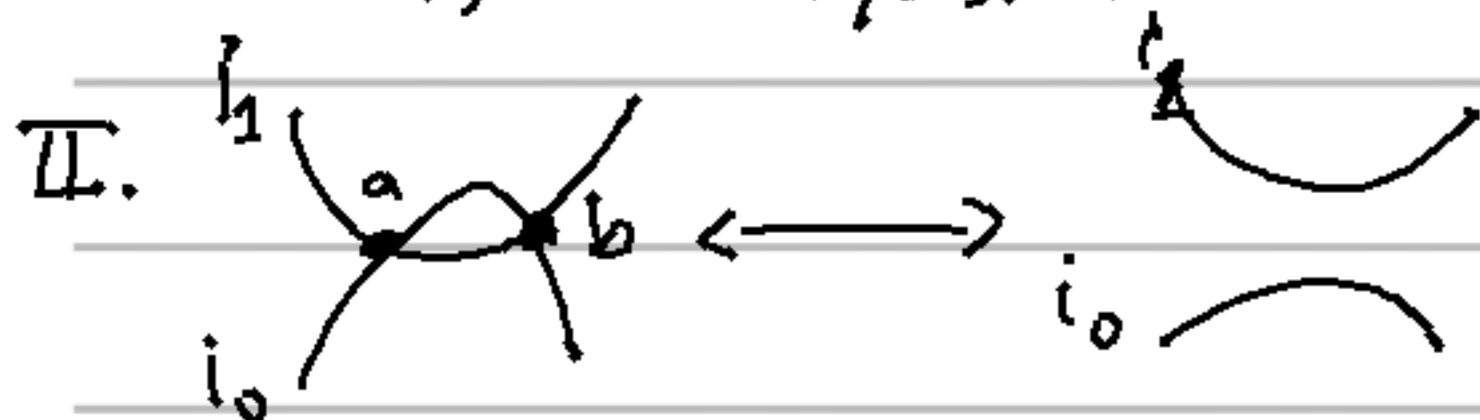
\Rightarrow dom map $\Rightarrow A_{\infty}$ morphism.

$\Rightarrow f, g$ are strict inverse A_{∞} isos.

For $x = a$, same set of argument

This case: strict iso.

For more \mathbb{I} , we need a quasi-iso.



$$\delta_a = b + v$$

Define $f: a \mapsto 0$

$b \mapsto v'$

$x \mapsto x'$ all other x .

Let $h: T_1 \rightarrow T_1[1]:$

$$h_1(b) = a$$

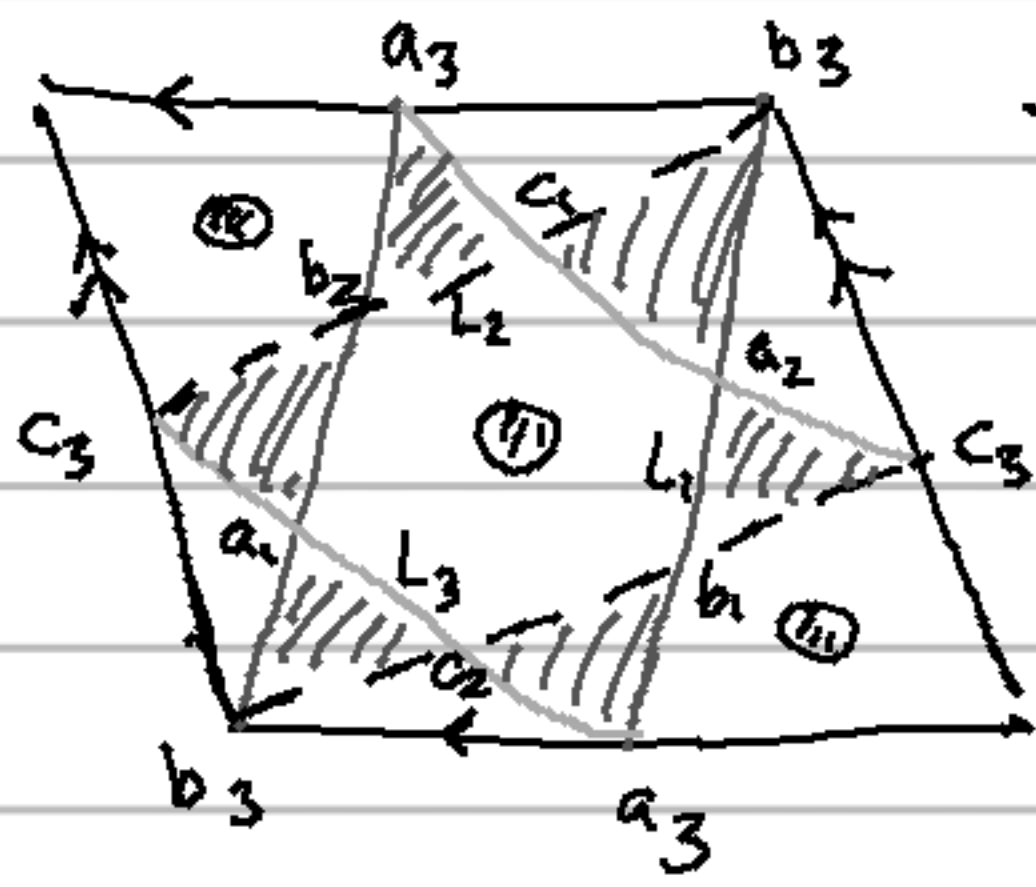
$$h_1(x) = 0 \text{ else } g_1: T_1' \rightarrow T_1$$

$$g_1(x') = x + h_1 \delta_1 x.$$

f is an A_{∞} -morphism, $f \circ g_1 = \text{id}$, $g_1 \circ f_1 = \text{id} + h \delta_1 + \delta_1 h_1 \Rightarrow$ By Perturbation Lemma, can extend g to $g: T' \rightarrow T$ st.

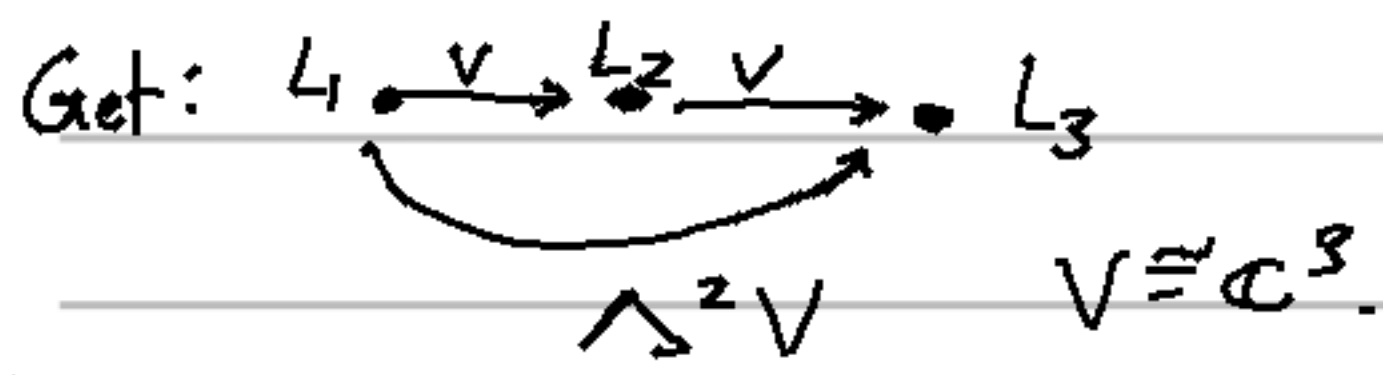
$f \circ g, g \circ f$ homotopic to id.

Quick Example:



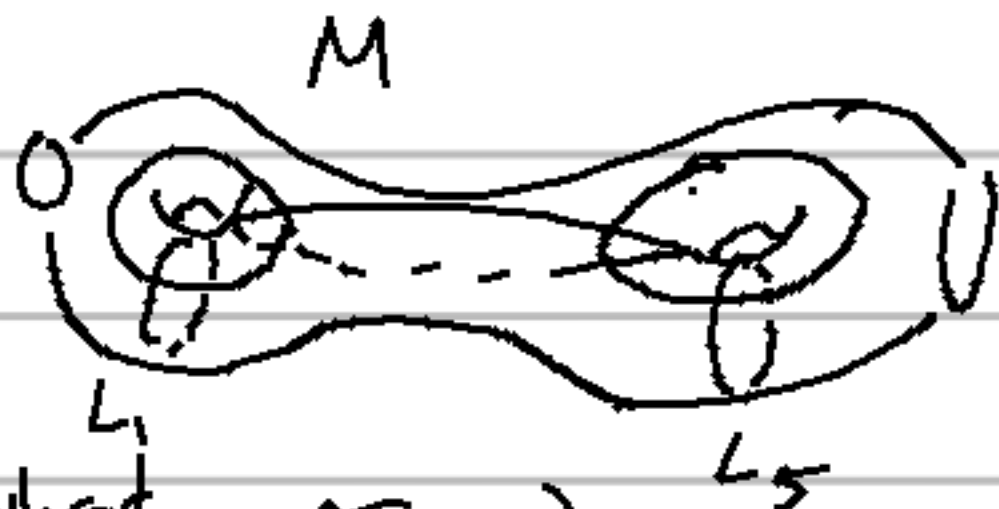
Torus w/ 3 points removed.

$$a_i b_j = \epsilon_{ijk} C_k$$



From Paul: Warning! Just b/c $F \rightarrow$ is combinatorial doesn't mean F is!

Ex:



full subcat
w/ obj. $\subset \mathcal{F}(M)$

$L_1 \rightarrow L_2$ Thm: This is not formal.

(even though there are no polygons in picture above. when you perturb, get small polygons)